AAE 526: More Linear Programming

Thomas F. Rutherford

Fall Semester, 2019 Department of Agricultural and Applied Economics University of Wisconsin, Madison

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• Another LP example

- Linear algebra and indexed GAMS code
- Linear programs in standard form
- Supersize me!
- Visualization



The Wyndor Glass Co. produces high-quality glass products, including windows and glass doors. It has three plants. Aluminum frames and hardware are made in Plant 1, wood frames are made in Plant 2 and Plant 3 cuts the glass and assembles the products.

Profits are 3 for product 1 and 5 per unit of product 2.

Because of declining earnings, top management has decided to revamp the product line. Unprofitable products are to be discontinued, and product capacity will be reassigned to launch two new products.

Describe Constraints

The data indicates that each batch of product 1 produced per week uses 1 hour of production time per week in Plant 1, whereas only 4 hours per week are available. The restriction is written mathematically as

$$x_1 \leq 4$$

Similarly, Plant 2 imposes the restriction

$$2x_2 \leq 12$$

The number of hours of production time used per week in Plant 3 given x_1 and x_2 is $3x_1 + 2x_2$ (3 hours of production per batch of good 1 and 2 hours per batch of good 2). The available capacity in Plant 3 provides the constraint:

$$3x_1 + 2x_2 \le 18$$

Finally, production rates cannot be negative, so it is necessary to impose restrictions $x_1 \ge 0$ and $x_2 \ge 0$





$\max Z = 3x_1 + 5x_2$

subject to the restrictions

and

$$x_1 \ge 0, x_2 \ge 0$$

gamside: D:\newproject.gpr - [d:\Optimization\lectures\Lecture2\wyndor.gms]							
📕 File	File Edit Search Windows Utilities Model Libraries Help						
>		A (a) 🚭 🕒					
wyndor.gr	ns						
VAR	IABLES						
	Z	Objective function (\$1000 per week)					
	X1	Glass doors (batches per week)					
	X2	Wood framed windows (batches per week);					
EQU	ATIONS						
	Zdef	Defines the objective function					
	plant1	Constraint imposed by plant 1					
	plant2	Constraint imposed by plant 2					
	plant3	Constraint imposed by plant 3;					
zde	f	Z =E= 3*X1 + 5*X2;					
pla	nt1	X1 =L= 4;					
pla	nt2	2*X2 =L= 12;					
pla	nt3	3*X1 + 2*X2 =L= 18;					
MOD	EL wyndor	<pre>/zdef, plant1, plant2, plant3/;</pre>					
SOL	SOLVE wyndor USING LP MAXIMIZING Z;						



This problem has only two dimensions, so a graphical procedure can be employed. We use label the axes as x_1 and x_2 . The first step is then to identify on the graph values of (x_1, x_2) which are *feasible* (consistent with the restrictions).



Values of (x_1, x_2) consistent with the constraints $0 \le x_1 \le 4$ and $0 \le x_2$:



The Feasible Region





Solution



The optimal solution is $x_1^* = 2, x_2^* = 6$ with $Z^* = 36$. This implies 2 batches of product 1 and 6 batches of product 2 will be produced per week, providing a total profit of \$36,000 per week.





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```
$TITLE Brewery Profit Maximization
       i
              Products /lager, ale/
set
               Ingredients /
       i.
                      malt
                            Malt,
                      yeast Yeast,
                      dehops German hops,
                      wihops Wisconsin hops/;
parameter
       p(j) Profit by product /lager 12, ale 9/
       s(i)
               Supply by ingredient /malt 4800, yeast 1750,
                                      dehops 1000, wihops 1750/;
table
               a(i,j) Requirements
                      lager
                              ale
       malt
                      4
                              2
                      1
                              1
       veast
                      1
       dehops
                              0
       wihops
                      0
                              1:
```

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```
Y(j) Production levels,
variables
              7.
                    Profit (maximand):
nonnegative variable Y;
equations
             supply(i) Ingredient supply
              profit Defines Z;
supply(i).. sum(j, a(i,j)*Y(j)) =L= s(i);
profit.. Z =E= sum(j, p(j)*Y(j));
MODEL BREWERY / supply, profit/;
solve BREWERY using LP maximizing Z;
```



A matrix is an array of numbers. $A \in \mathbb{R}^{m \times n}$.

$$A = \begin{bmatrix} a_{11} & \cdots & a_{1n} \\ \vdots & \ddots & \vdots \\ a_{m1} & \cdots & a_{mn} \end{bmatrix}$$

which has *m* rows and *n* columns.

The table statement in GAMS can be used to define a matrix:

Table versus Parameter



A matrix may be specified either the table or parameter statement:

```
set i Row indices /1*3/,
      j Column indice /a,b/;
table a(i,j) Matrix with three rows and two columns
                b
           а
      1 0.23 12.3
      2 -0.1 2.4
      3 3.2 0.1;
parameter b(i,j) The same matrix in database format /
      1.a 0.23
      2.a -0.1
      3.a 3.2
      1.b 12.3
      2.b 2.4
      3.b 0.1 /;
parameter c(i,j) Check that a=b; c(i,j) = a(i,j) - b(i,j); display c;
       22 PARAMETER c check that a=b
                 (ALL 0.000)
```



Two matrices can be multiplied *if their inner dimensions agree*. In matrix notation (MATLAB style):



In detached coefficient notation (GAMS style) we write:

$$c_{ij} = \sum_k a_{ik} b_{kj}$$

In GAMS syntax, we have:

```
c(i,j) = sum(k, a(i,k) * b(k,j));
```



- The transpose operator A^T swaps rows and columns. If $A \in \mathbb{R}^{m \times n}$, then $A^T \in \mathbb{R}^{n \times m}$ and $(A^T)_{ij} = A_{ji}$
- It follows that:

$$(A^{T})^{T} = A$$
$$(AB)^{T} = B^{T}A^{T}$$

Linear and affine functions

• A function $f(x_1, ..., x_m)$ is *linear* in the variables $x_1, ..., x_m$ if there exists constants $a_1, ..., a_m$ such that

$$f(x_1,\ldots,x_m)=a_1x_1+\ldots+a_mx_m=\sum_i a_ix_i=a^Tx$$

• A function $f(x_1, ..., x_m)$ is *affine* in the variables $x_1, ..., x_m$ if there exists constants $b, a_1, ..., a_m$ such that

$$f(x_1,...,x_m) = ba_1x_1 + ... + a_mx_m = b + \sum_i a_ix_i = b + a^Tx_i$$

- Examples:
 - 3x y is *linear* in (x, y).
 2xy + 1 is *affine* in x and y, but not in (x, y).
 x² + y² is neither linear nor affine.





Several linear or affine functions can be combined:

 $\begin{array}{cccc} a_{11}x_1 & + \dots & +a_{1n}x_n + b_1 \\ a_{21}x_1 & + \dots & +a_{2n}x_n + b_2 \\ \vdots & \vdots & \\ a_{m1}x_1 & + \dots & +a_{mn}x_n + b_m \end{array} \Rightarrow \left[\begin{array}{c} a_{11} & \cdots & a_{1n} \\ \vdots & \ddots & \vdots \\ a_{m1} & \cdots & a_{mn} \end{array} \right] \left[\begin{array}{c} x_1 \\ \vdots \\ x_n \end{array} \right] + \left[\begin{array}{c} b_1 \\ \vdots \\ b_n \end{array} \right]$

which can be written simply as Ax + b. Same definitions apply:

- A vector-valued function F(x) is linear in x if there exists a constant matrix A such that F(x) = Ax.
- A vector-valued function F(x) is affine in x if there exists a constant matrix A and vector b such that F(x) = Ax + b.

W

Matrix basics: inner and outer products

A vector is a column matrix. We write $x \in \mathbb{R}^n$ to mean that

$$\mathbf{x} = \left[\begin{array}{c} \mathbf{x}_1 \\ \vdots \\ \mathbf{x}_n \end{array} \right]$$

This is an $n \times 1$ matrix.

Two vectors $x, y \in \mathbb{R}^n$ can be multiplied together in two ways. Both are valid matrix multiplications:

- inner product: produces a scalar, $x^T y = x_1 y_1 + \cdots + x_n y_n$.
- **outer product**: produces an *n* × *n* matrix.

$$xy^{T} = \begin{bmatrix} x_{1}y_{1} & \cdots & x_{1}y_{n} \\ \vdots & \ddots & \vdots \\ x_{n}y_{1} & \cdots & x_{n}y_{n} \end{bmatrix}$$

Calculating Inner and Outer Products in GAMS

```
set i/i1*i3/:
parameter x(i), y(i), xy, xyt(i,i);
       Generate two random arrays containing values
*
       between zero and one:
*
x(i) = uniform(0,1); y(i) = uniform(0,1);
       Inner product:
*
xy = sum(i, x(i)*y(i));
       Need a second symbol to refer to the set i:
*
alias (i,j);
* Outer product:
xyt(i,j) = x(i)*y(j);
display x, y, xy, xyt;
```

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	23 PARA	METER x				
i1 0.172	, i2	0.843, i3	0.550			
	23 PARA	METER y				
i1 0.301	, i2	0.292, i3	0.224			
	23 PARA	METER xy		=	0.421	
	23 PARA	METER xyt				
	i1	i2	i3			
i1	0.052	0.050	0.038			
i2	0.254	0.246	0.189			
i3	0.166	0.161	0.123			





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The linear program

A linear program is an optimization model with:

- real-valued variables ($x \in \mathbf{R}^n$)
- linear cost function $(c^T x)$
- constraints may be:
 - affine equations (Ax = b)
 - affine inequalities $(Ax \le b \text{ or } Ax \ge b)$
 - combinations of the above
- individual variables may have:
 - bounds $(p \leq x_i$, or $x_i \leq q$, or $p \leq x_i \leq q)$
 - no bounds (x_i is unconstrained)

There are many equivalent representations of any linear program.



Linear Programming



Standard Form:

Why it's hard:

- Lots of variables (*n* of them)
- Lots of boundaries to check (the inequalities)

Why it's not impossible:

• All expressions are linear

For any given linear programming problem, exactly one of the following statements applies:

- 1. The model is infeasible: there is no x that satisfies all the constraints. (is the model correct?)
- The model is feasible, but unbounded: the cost function can be arbitrarily improved. (forgot a constraint?)
- Model has a solution which occurs on the boundary of the feasible polyhdron. Note that there is no guarantee that the solution is unique – there may be many solutions!



infeasible



unbounded



boundary





• Every linear program can be put into the form:

$$\max_{z \in \mathbb{R}^n} c^T z$$

subject to:

$$Az \leq b$$

 $z > 0$

• This is call the *standard form* of a linear program.

Brewery Profit: Standard Form



$$\begin{array}{rll} \max_{x,y} & 120x + 90y \\ {\rm s.t.:} & 4x + 2y & \leq & 4800 \\ & x + y & \leq & 1750 \end{array}$$

 $0 \le x \le 1000, \quad 0 \le y \le 1500$

is equivalent to:

$$\max_{x,y} \begin{bmatrix} 120 \\ 90 \end{bmatrix}^{T} \begin{bmatrix} x \\ y \end{bmatrix}$$

s.t.
$$\begin{bmatrix} 4 & 2 \\ 1 & 1 \\ 1 & 0 \\ 0 & 1 \end{bmatrix} \begin{bmatrix} x \\ y \end{bmatrix} \leq \begin{bmatrix} 4800 \\ 1750 \\ 1000 \\ 1500 \end{bmatrix}$$
$$x, y \ge 0$$

Hence, our brewery profit maximization model can be transformed into standard inequality form with the assignments:

$$z = \begin{bmatrix} x \\ y \end{bmatrix} \qquad c = \begin{bmatrix} 120 \\ 90 \end{bmatrix}$$
$$A = \begin{bmatrix} 4 & 2 \\ 1 & 1 \\ 1 & 0 \\ 0 & 1 \end{bmatrix} \qquad b = \begin{bmatrix} 4800 \\ 1750 \\ 1000 \\ 1500 \end{bmatrix}$$

A partial list, taken from Ferris et al., Chapter 1:

- Resoure allocation
- The diet problem
- Linear surface fitting
- Load balancing
- Classification

. . .

• Minimum cost network flow





A company has *m* products which are produced with *n* resources. The value of product *i* is c_i , while each unit of resource *j* costs d_j dollars. One unit of product *i* requires a_{ij} units of resource *j*, and a maximum of b_j units of resource *j* are available:

$$\max_{x,y} z = \sum_{i} c_{i} y_{i} - \sum_{j} d_{j} x_{j}$$

subject to

$$x_j = \sum_i a_{ij} y_i, \quad x_j \leq b_j, \quad x_j \geq 0, y_i \geq 0$$

Note that the constraints can be written in *detached coefficient form* as:

$$x_j = \sum_i a_{ij}y_i = a_{1j}y_1 + a_{2j}y_2 + \ldots + a_{mj}y_m$$



Given the prices p_j of food type j, the content of nutrient i in food $j(a_{ij})$ and the dietary requirement of nutrient i, b_i , solve:

$$\min_{x} z = \sum_{j} p_{j} x_{j}$$

subject to

$$\sum_{j} a_{ij} x_j \ge b_i, \quad \forall i$$
$$x_j \ge 0$$

Linear Surface Fitting



Given a set of observations $A = [a_{ij}]$ and b_i . Find weights on the columns of A and a scalar constant γ which best "predicts" the value of b on the basis of observations a_{ij} , assuming a linear model:

$$\min_{\mathbf{x},\gamma}\sum_{i=1}^{m}\left|\sum_{j}a_{ij}x_{j}+\gamma-b_{i}\right|$$

or, equivalently

$$\min_{x,\gamma,y} z = \sum_i y_i$$

subject to

$$-y_i \leq \sum_j a_{ij}x_j + \gamma - b_i \leq y_i$$

Note that the constraint ensures that each y_i is no smaller than the absolute value $|\sum_j a_{ij}x_j + \gamma - b_i|$.

Load Balancing



Balance computational work among n processors, distributing the load in such a way that the lightest-loaded processor has as heavy a load as possible:

- p_i Current load of processor $i = 1, 2, \ldots, n$
- *L* Total load to be distributed
- x_i Fraction of additional load L to be distributed to processor i, with $x_i \ge 0$ and $\sum_i x_i = 1$.

 $\gamma\,$ minimum final loads after distribution of the new workload L

 $\max_{x,\gamma}\gamma$

subject to

$$\gamma \leq p_i + x_i L, \quad \sum_i x_i = 1, \quad x_i \geq 0 \ \forall i$$



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In words:

Minimize: the cost (or calories) of eating at McDonald's

Subject to: the total amounts of food or nutrients fall between certain minimum and maximum values



A GAMS Model



set	n	Nutritional needs / calo Calories carbo Carbohydrates protein Protein vita Vitamin A vitc Vitamin C calc Calcium iron Iron /,
	f	<pre>Foods / "Quarter Pounder w/ Cheese" "McLean Deluxe w/ Cheese" "Big Mac", "Filet-O-Fish", "McGrilled Chicken", "Fries, small", "Sausage McMuffin", "1% Lowfat Milk", "Orange Juice" /;</pre>



http://nutrition.mcdonalds.com/getnutrition/nutritionfacts.pdf



McDonald's USA Nutrition Facts for Popular Menu Items

We provide a nutrition analysis of our menu items to help you balance your McDonald's meal with other foods you eat. Our goal is to provide you with the information you need to make sensible decisions about balance, variety and moderation in your diet.

			Fat		1	(B)	1		(Bu	ŧ,		1	(B) si	1	(6)	ŧ.			%	DAIL	VAL	JE
Nutrition Facts	Serving Size	Calories	Calories from	Total Fat (g)	% Daily Value	Saturated Fat	% Daily Value	Trans Fat (g)	Cholesterol (% Daily Value	Sodium (mg)	% Daily Value	Carbohydrate	% Daily Value	Dietary Fiber	% Daily Value	Sugars (g)	Protein (g)	Vitamin A	Vitamin C	Calcium	Iron
Burgers & Sandy	viches																					
Angus Chipotle BBQ Bacon†+++	10.3 oz (294 g)	800	350	39	60	18	88	2	145	49	2020	84	66	22	4	14	16	45	10	2	25	35

The Data

W

table a(f,n) Nutritional content of foods

	calo	carbo	protein	vita	vitc	calc	iron
"Quarter Pounder w/ Cheese"	510	34	28	15	6	30	20
"McLean Deluxe w/ Cheese"	370	35	24	15	10	20	20
"Big Mac"	500	42	25	6	2	25	20
"Filet-O-Fish"	370	38	14	2	0	15	10
"McGrilled Chicken"	400	42	31	8	15	15	8
"Fries, small"	220	26	3	0	15	0	2
"Sausage McMuffin"	345	27	15	4	0	20	15
"1% Lowfat Milk"	110	12	9	10	4	30	0
"Orange Juice"	80	20	1	2	120	2	2;

table nutr(n,*) Nutrient requirements

	min	шах
calo	2000	inf
carbo	350	375
protein	55	inf
vita	100	inf
vitc	100	inf
calc	100	inf
iron	100	inf;

table food(f,*) Food cost and requirements

	cost	min	max
"Quarter Pounder w/ Cheese"	1.84	0	inf
"McLean Deluxe w/ Cheese"	2.19	0	inf
"Big Mac"	1.84	0	inf
"Filet-O-Fish"	1.44	0	inf
"McGrilled Chicken"	2.29	0	inf
"Fries, small"	0.77	0	inf
"Sausage McMuffin"	1.29	0	inf
"1% Lowfat Milk"	0.6	0	inf
"Orange Juice"	0.72	0	inf;

Model Equations



nonnegative variables	Y(n) X(f)	Nutritional content Purchased quantity;			
free variable	COST	Total cost;			
equations	ydef, objdef;				
ydef(n)	Y(n) =e=	= sum(f, X(f)*a(f,n));			
objdef	COST =e=	<pre>sum(f, X(f) * food(f,"cost"));</pre>			
<pre>model mincost /ydef, objdef /;</pre>					
Y.LO(n) = nutr(n,"min"); Y.UP(n) = nutr(n,"max"); X.LO(f) = food(f,"min"); X.UP(f) = food(f,"max");					



solve mincost using lp minimizing COST;

* Generate reports of the menu and diet:

```
parameter menu Resulting menu,;
menu("Cost","MinCost") = COST.L;
menu("Calories","MinCost") = Y.L("calo");
menu("ModelStat","MinCost") = mincost.modelstat;
menu("solvestat","MinCost") = mincost.solvestat;
menu(f,"MinCost") = X.L(f);
display menu;
```



	MinCost
Cost	14.856
Calories	3965.369
Quarter Pounder w/ Cheese	4.385
Fries, small	6.148
1% Lowfat Milk	3.422

Cheap, but 4000 calories!



Put an upper bound on calories.

```
Y.UP("calo") = 2500;
solve mincost using lp minimizing COST;
```

Calories are down, cost is up and the diet looks better.

	MinCost	Cal2500
Cost	14.856	16.671
Calories	3965.369	2500.000
Quarter Pounder w/ Cheese	4.385	0.232
McLean Deluxe w/ Cheese		3.855
Fries, small	6.148	
1% Lowfat Milk	3.422	2.043
Orange Juice		9.134

Third Run



Try for a 2000 calorie diet.

```
Y.UP("calo") = 2000;
solve mincost using lp minimizing COST;
```

Not possible!

MODEL	mincost	OBJECTIVE	С
TYPE	LP	DIRECTION	MINIMIZE
SOLVER	CPLEX	FROM LINE	121

****	SOLVER STATUS	1 Normal Completion
****	MODEL STATUS	4 Infeasible
****	OBJECTIVE VALUE	0.9121

Minimize calories, ignoring cost.



variable	CALO	Objecti	ve value	calories;
equation	objcalo	Objecti	ve minimi	ze calories;
objcalo	CALO =e= Y("	calo");		
model mincal /yo	lef, objdef,	objcalo/;		
solve mincal us:	ing lp minimi	zing CALO;		
Minimum calories is	s 2467 at a cost	of \$16.75:		
		MinCost	Cal2500	MinCal
Cost		14.856	16.671	16.745
Calories		3965.369	2500.000	2466.981
Quarter Pounder	w/ Cheese	4.385	0.232	
McLean Deluxe w	/ Cheese		3.855	4.088
Fries, small		6.148		
1% Lowfat Milk		3.422	2.043	2.044
Orange Juice			9.134	9.119



X.UP(f) = 2; solve mincost using lp minimizing COST;

More intersting cuisine. Cost is up a bit, calories are down relative to original solution.

	MinCost	MinCal	MinCostV
Cost	14.856	16.745	16.766
Calories	3965.369	2466.981	3798.077
Quarter Pounder w/ Cheese	4.385		2.000
McLean Deluxe w/ Cheese		4.088	2.000
Big Mac			2.000
Fries, small	6.148		1.423
Sausage McMuffin			1.000
1% Lowfat Milk	3.422	2.044	2.000
Orange Juice		9.119	2.000



X.UP(f) = 2; solve mincal using lp minimizing COST;

Almost 3500 calories. Wow!

	MinCost	MinCal	MinCostV	MinCalV
Cost	14.856	16.745	16.766	17.248
Calories	3965.369	2466.981	3798.077	3488.287
Quarter Pounder w/ Cheese	4.385		2.000	1.952
McLean Deluxe w/ Cheese		4.088	2.000	2.000
Big Mac			2.000	
Filet-O-Fish				0.359
McGrilled Chicken				2.000
Fries, small	6.148		1.423	2.000
Sausage McMuffin			1.000	
1% Lowfat Milk	3.422	2.044	2.000	2.000
Orange Juice		9.119	2.000	2.000

Whole Numbers Solution



Move from linear program (LP) to mixed integer programming (MIP).

```
integer variable XI(f) Food purchase;
equation xidef; xidef(f).. X(f) === XI(f);
```

```
XI.LO(f) = 0; XI.UP(f) = 2;
```

model integerdiet /ydef, objdef, objcalo, xidef /; solve integerdiet using MIP minimizing CALO;

This only makes a small increase in price and calories.

	MinCost	MinCal	MinCostV	MinCalV	MinCalInt
Cost	14.856	16.745	16.766	17.248	17.490
ModelStat	1.000	1.000	1.000	1.000	1.000
SolveStat	1.000	1.000	1.000	1.000	1.000
Quarter Pounder w/ Cheese	4.385		2.000	1.952	2.000
McLean Deluxe w/ Cheese		4.088	2.000	2.000	2.000
Big Mac			2.000		
Filet-O-Fish				0.359	1.000
McGrilled Chicken				2.000	2.000
Fries, small	6.148		1.423	2.000	1.000
Sausage McMuffin			1.000		
1% Lowfat Milk	3.422	2.044	2.000	2.000	2.000
Orange Juice		9.119	2.000	2.000	2.000
Calories	3965.369	2466.981	3798.077	3488.287	3530.000

Ŵ

Include a set definition to sequence rows in the report – the *universal element list*:

<pre>set seq /Cost, Calo, ModelStat, SolveStat/;</pre>				
	MinCost	Cal2500	Cal2000	
Cost	14.856	16.671	17.796	
Calo	3965.369	2500.000	2000.000	
ModelStat	1.000	1.000	4.000	
SolveStat	1.000	1.000	1.000	

Some GAMS Program Details



Uses a *batinclude subroutine* to add scenario results to the report parameters menu and diet:

```
parameter menu Resulting menu,
                diet Resulting diet;
$onechov >%gams.scrdir%report.gms
menu("Cost", "%1") = C.L;
menu("Calories","%1") = Y.L("calo");
menu("ModelStat","%1") = mincost.modelstat;
menu("solvestat","%1") = mincost.solvestat;
menu(f, "%1") = X.L(f);
diet("Cost","%1") = C.L:
diet("ModelStat","%1") = mincost.modelstat;
diet("SolveStat", "%1") = mincost.solvestat;
diet(n, "%1") = Y.L(n);
$offecho
        Define an environment variable to label the report:
solve mincost using lp minimizing COST;
$batinclude %gams.scrdir%report MinCost
```

W

Compute the integer programming solution by introducing an integer variable and assigning it to the LP decision variable. All variables and equations from the LP remain in the model so reporting routine is unchanged!

```
integer variable XI(f) Food purchase;
equation xidef;
xidef(f).. X(f) === XI(f);
XI.LO(f) = 0;
XI.UP(f) = 2;
model integerdiet /ydef, objdef, objcalo, xidef /;
solve integerdiet using MIP minimizing CAL;
```



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subject to:

 $\max_{x,y} 120x + 90y$ $4x + 2y \le 4800$ $x + y \le 1750$ $0 \le x \le 1000$ $0 \le y \le 1500$

In which:

- x : number of batches of lager produced
- y : number of batches of ales produced

Visualization: Scatter Plots in Excel





Visualization: GNUPLOT





X axis

3d Hyperplane (alpha= 0.20, beta= 0.50)



Z axis

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Visualization: GNUPLOT





GNUPLOT Command File



```
#
        Begin with a reset so we can make changes and immediately reload:
reset
#
        Define parameters are user-defined GNUPLOT variables:
maltg=4800
veastg=1750
maltx=4
maltv=2
profitx=12
profity=9
        Calculate the optimum, assuming that it is where the yeast
#
        and malt constraints are both binding:
#
xmax=(maltg-malty*veastg)/(maltx-malty)
vmax=veastg-xmax
maxprofit=xmax*profitx+ymax*profity
        Calculate points where the hops constraints intersect
#
        the yeast and malt constraints:
#
xlim = veastg-1500
ylim = (maltq-maltx*1000)/malty
set style arrow 1 nohead linecolor rgb "gray" linewidth 2 dashtype solid
set arrow 1 from 1000,0,0 to 1000,ylim,0 arrowstyle 1
set arrow 2 from 0.1500.0 to xlim.1500.0 arrowstyle 1
```

GNUPLOT Command File (cont)

Define linear functions representing the yeast and malt
constraints as well as the isoprofit line at the optimal
point:
yeast(x)=yeastq-x
malt(x)=(maltq-maltx*x)/malty
isoprofit(x) = ymax-profitx*(x-xmax)/profity

Set up axes:

```
set xrange [0:1500]
set yrange [0:2500]
set ylabel 'batches of ale (y)' offset -1,0
set xlabel 'batches lager (x)' offset 0,-1
set xtics axis
set ytics axis
unset border
set arrow 3 from 0,0 to 1450,0 linestyle 1
set arrow 4 from 0,0 to 0,2550 linestyle 1
```

Label the feasible region:

set label "feasible region" at 200,500

Put a black circle at the optimum:

```
optimum = sprintf('(%.f,%.f)',xmax,ymax)
set label optimum at xmax+30,ymax+30
set object circle at xmax,ymax front size 6 \
    fillstyle solid 1 fillcolor rgb "black"
```



```
W
```

```
#
        Generate a data file with extreme points of the feasible region:
set print "feasible.dat"
print sprintf('%.f %.f',0,0)
print sprintf('%.f %.f',0,1500)
print sprintf('%.f %.f',xlim, 1500)
print sprintf('%.f %.f'.xmax, vmax)
print sprintf('%.f %.f',1000, vlim)
print sprintf('%.f %.f',1000, 0)
unset print
        Define the style to be used for "filledcurves" to denote the
#
        feasible region:
#
set style fill transparent solid 0.2 noborder
set style line 1 linecolor "black" linewidth 2 dashtype 1
set style line 2 linecolor "forest-green" linewidth 2 dashtype 1
set style line 3 linecolor "red" linewidth 2 dashtype 2
plot yeast(x) ls 1, malt(x) ls 2, isoprofit(x) ls 3, \setminus
      'feasible.dat' using 1:2 with filledcurves below notitle
```